



Universidade do Minho  
Escola de Engenharia

# Noise Resilient Quantum Amplitude Estimation

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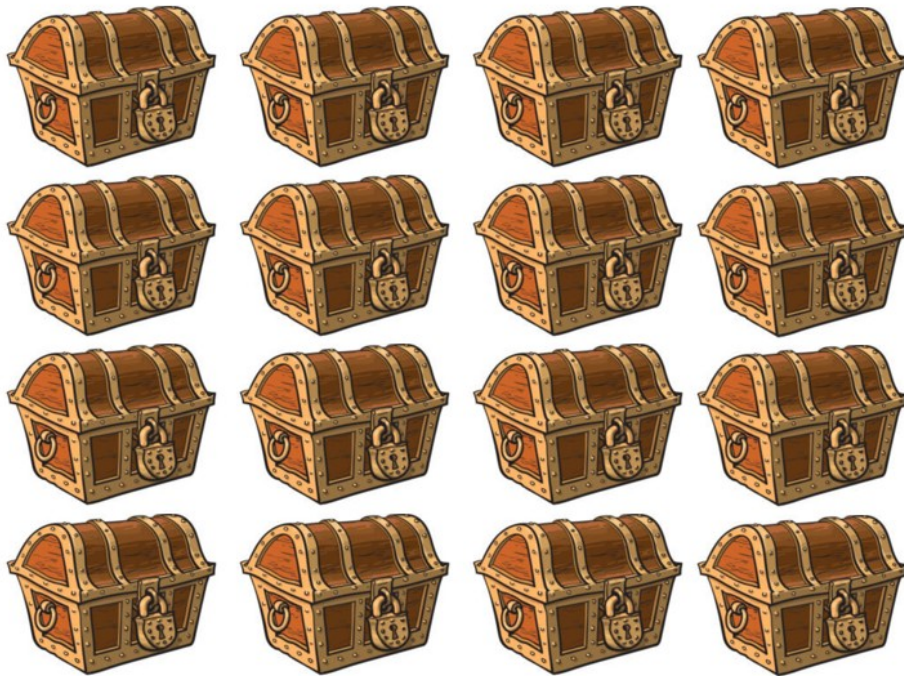
**Work supervised by:**

- *Luís P. Santos*
- *Ernesto F. Galvão*

# Overview

1. Quantum searching
2. Quantum amplitude estimation
3. Noisy quantum devices
4. Numerical results
5. Overview and numerical analysis of QAE algorithms
6. Bayesian amplitude estimation

# Quantum Searching



Quantum resources can bring a **quadratic speed up** to the task of **searching an unstructured**

# Quantum Searching

Mathematically:

$$f(x) : \{0, 1\}^n \rightarrow \{0, 1\}$$

$$f(x) = \begin{cases} 1, & \text{if } x \text{ is a solution;} \\ 0, & \text{otherwise.} \end{cases}$$



;



find  $x$  s. t.  $f(x) = 1$



# Quantum Searching

$$f(x) = \begin{cases} 1, & \text{if } x \text{ is a solution;} \\ 0, & \text{otherwise.} \end{cases}$$

## Classical approach:

1. Sample uniformly at random.
2. If , terminate.
3. Go to 1.

$$\begin{array}{l} \text{Cost} \\ : \\ \sum_{k=1}^N \frac{1}{N} \cdot k = \frac{1}{N} \cdot \left( \frac{N(N+1)}{2} \right) = \frac{N+1}{2} \end{array}$$

$$\in O(N)$$

or if there are >1 solutions:

$$O(N/M)$$

# Quantum Searching

Assume we can encode  $f$  as a quantum phase oracle .

$$f(x) = \begin{cases} 1, & \text{if } x \text{ is a solution;} \\ 0, & \text{otherwise.} \end{cases} \quad \hat{U}_f |x\rangle = (-1)^{f(x)} |x\rangle$$

# Quantum Searching

**Quantum approach:**

Prepare a superposition  $|\psi_A\rangle = A|0\rangle^{\otimes n}$  containing the solution.

# Quantum Searching

Quantum approach:

Prepare a superposition  $|\psi_A\rangle = A|0\rangle^{\otimes n}$  containing the solution.

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E.g. use the Hadamard transform:  $|\psi_A\rangle = H^{\otimes n}|0\rangle^{\otimes n} = \frac{1}{\sqrt{N}} \sum_{x \in \{0,1\}^n} |x\rangle$

Call the subset of solutions, its complementary in . Define:

$$|\Psi_1\rangle = \sqrt{\frac{1}{M}} \sum_{x \in X} |x\rangle \quad ; \quad |\Psi_0\rangle = \sqrt{\frac{1}{N-M}} \sum_{x \in X^c} |x\rangle$$

Then

:

$$|\psi_A\rangle = \sqrt{a} |\psi_1\rangle + \sqrt{1-a} |\psi_0\rangle$$

$$a \equiv \frac{M}{N}$$



# Quantum Searching

Quantum approach:

Prepare a superposition

$$|\psi_A\rangle = A|0\rangle^{\otimes n}$$

containing the solution.

$$\begin{aligned}
 & \frac{1}{\sqrt{N}} \sum_{c=0}^{N-1} |c\rangle \\
 &= \frac{1}{\sqrt{N}} \left( \sqrt{a} |c_{\text{solution}}\rangle + \sqrt{1-a} |c_{\text{not solution}}\rangle \right) \\
 &= \frac{1}{\sqrt{N}} \left( \sqrt{a} \left| \begin{array}{c} 0 \\ \vdots \\ 0 \\ 1 \\ 0 \\ \vdots \\ 0 \end{array} \right\rangle + \sqrt{1-a} \left| \begin{array}{c} 0 \\ \vdots \\ 0 \\ 0 \\ 1 \\ \vdots \\ 0 \end{array} \right\rangle \right)
 \end{aligned}$$

The diagram illustrates the preparation of a superposition state. It shows a sum over all possible bit strings  $|c\rangle$  from  $0$  to  $N-1$ . The state is then decomposed into two parts: a component with amplitude  $\sqrt{a}$  for the state where the solution bit is 1, and a component with amplitude  $\sqrt{1-a}$  for the state where the solution bit is 0. Treasure chest icons are used to highlight the solution bit in each state.

# Quantum Searching

**Quantum approach:**

Prepare a superposition  $|\psi_A\rangle = A|0\rangle^{\otimes n}$  containing the solution.

Amplify the pre-image of 1 under  $A$  via quantum amplitude amplification (Grover's algorithm).

# Quantum Searching

Quantum approach:

Prepare a superposition  $|\psi_A\rangle = A|0\rangle^{\otimes n}$  containing the solution.

Amplify the pre-image of 1 under  $U_f$  via quantum amplitude amplification (QAA)

Define  $\theta = \arcsin(\sqrt{a}) \leftrightarrow a = \sin^2(\theta)$

Then  $|\psi_A\rangle = \sqrt{a}|\psi_1\rangle + \sqrt{1-a}|\psi_0\rangle$

becomes

:

$$|\psi_A\rangle = \sin(\theta)|\psi_1\rangle + \cos(\theta)|\psi_0\rangle$$

$$\hat{G} = -A\hat{U}_0A^{-1}\hat{U}_f$$

$$\hat{G}^m|\psi_A\rangle = \sin((2m+1)\theta)|\psi_1\rangle + \cos((2m+1)\theta)|\psi_0\rangle$$

$$|\langle\psi_1|\hat{G}^m|\psi_A\rangle|^2 = \sin^2((2m+1)\theta)$$

$$(2m_{\text{ideal}} + 1)\theta = \pi/2$$

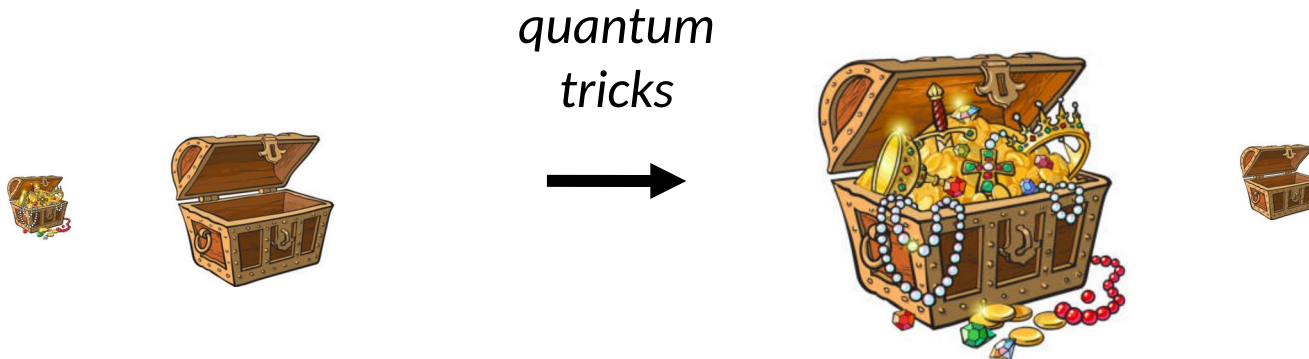
# Quantum Searching

Quantum approach:

Prepare a superposition  $|\psi_A\rangle = A|0\rangle^{\otimes n}$  containing the solution.

Amplify the pre-image of 1 under  $U$  via quantum amplitude amplification ( $U^2$ )

---



# Quantum Searching

Quantum approach:

Prepare a superposition  $|\psi_A\rangle = A|0\rangle^{\otimes n}$  containing the solution.

Amplify the pre-image of 1 under  $U$  via quantum amplitude amplification (QAA)

$$\begin{aligned}
 & \text{-----} \\
 & |\psi\rangle = \left| \frac{0}{c} \begin{array}{c} \text{treasure chest} \\ \hline \end{array} + \sqrt{1 - \frac{1}{c^2}} \begin{array}{c} \text{empty chest} \\ \hline \end{array} \right\rangle \\
 & \quad \quad \quad \downarrow \text{QAA} \\
 & |\psi\rangle = \sqrt{1 - a_{\text{amp}}} \begin{array}{c} \text{treasure chest} \\ \hline \end{array} + \sqrt{a_{\text{amp}}} \begin{array}{c} \text{empty chest} \\ \hline \end{array}
 \end{aligned}$$

# Quantum Searching

## Quantum approach:

Prepare a superposition  $|\psi_A\rangle = A|0\rangle^{\otimes n}$  containing the solution.

Amplify the pre-image of 1 under  $A$  via quantum amplitude amplification ( $A^{-1}$ ).

Measure to find a solution with high probability.

# Quantum Searching

Quantum approach:

Prepare a superposition  $|\psi_A\rangle = A|0\rangle^{\otimes n}$  containing the solution.

Amplify the pre-image of 1 under  $U$  via quantum amplitude amplification ( $U$  is the oracle).

Measure to find a solution with high probability.

$$|\psi\rangle = \sqrt{1 - a_{\text{amp}}} |0\rangle + \sqrt{a_{\text{amp}}} |1\rangle$$

to obtain  with probability.



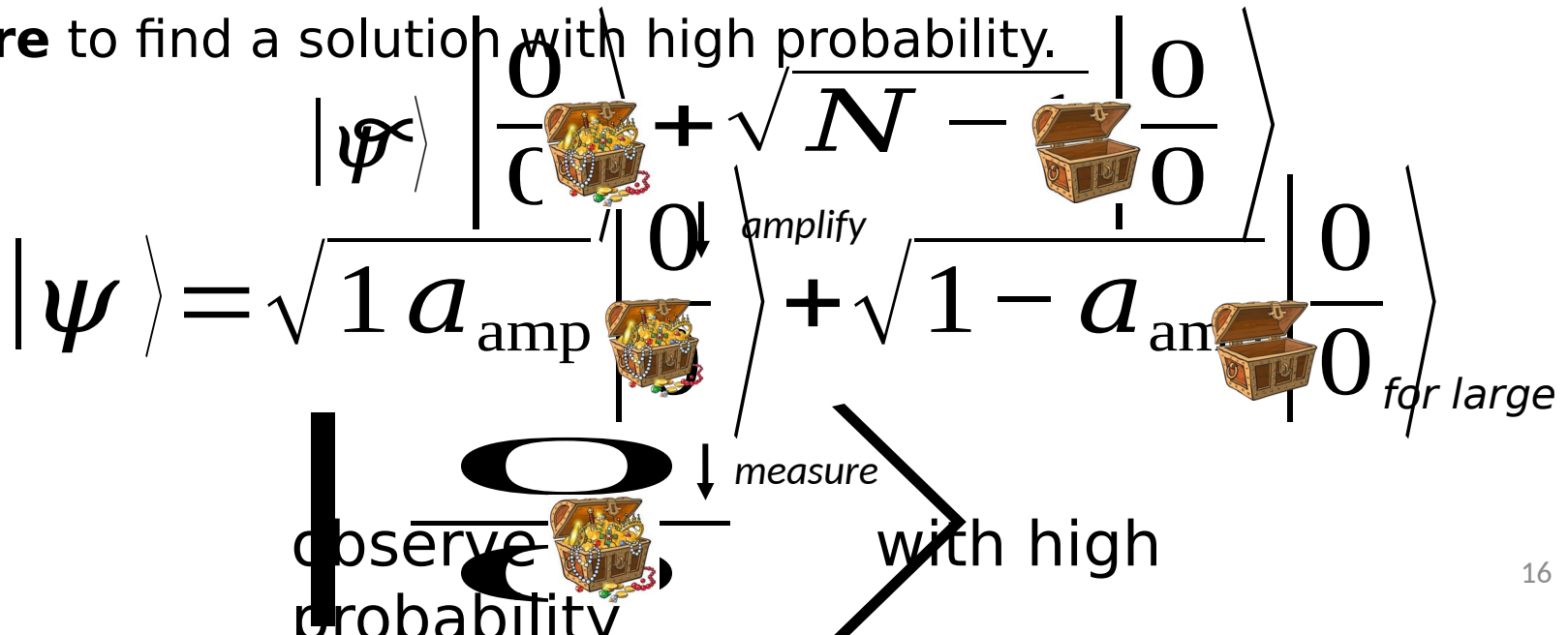
# Quantum Searching

Quantum approach:

Prepare a **superposition**  $|\psi_A\rangle = A|0\rangle^{\otimes n}$  containing the solution.

**Amplify** the pre-image of 1 under  $f$  via quantum amplitude amplification

**Measure** to find a solution with high probability.





# Quantum Searching

$$|\psi\rangle = \frac{1}{\sqrt{C}} \left| \begin{array}{c} 0 \\ \text{treasure chest} \end{array} \right\rangle + \sqrt{\frac{N-1}{N}} \left| \begin{array}{c} 0 \\ \text{empty chest} \end{array} \right\rangle$$

↓ amplify

$$|\psi\rangle = \sqrt{1 - a_{\text{amp}}} \left| \begin{array}{c} 0 \\ \text{treasure chest} \end{array} \right\rangle + \sqrt{1 - a_{\text{amp}}} \left| \begin{array}{c} 0 \\ \text{empty chest} \end{array} \right\rangle$$

*for large*

↓ measure

$$\left| \begin{array}{c} 0 \\ \text{treasure chest} \end{array} \right\rangle$$

observe with high probability

Cost  
:

$$\Theta(\sqrt{1/a})$$

$$\Theta(\sqrt{N/M})$$

# Quantum Amplitude Estimation

**Quantum amplitude estimation (QAE)** is related to quantum searching; it occurs in the same framework, and offers the same **quadratic speed-up**. It consists

on the task of estimating the parameter :

$$|\psi\rangle = \sqrt{a} \left| \begin{array}{c} 0 \\ \text{treasure chest} \end{array} \right\rangle + \sqrt{1-a} \left| \begin{array}{c} 0 \\ \text{empty chest} \end{array} \right\rangle$$

$$a = ?$$

# Quantum Amplitude Estimation

$$|\psi\rangle = \sqrt{a} \left| \begin{array}{c} 0 \\ \text{Treasure Chest} \end{array} \right\rangle + \sqrt{1-a} \left| \begin{array}{c} 0 \\ \text{Empty Chest} \end{array} \right\rangle \quad \boxed{a = ?}$$

It becomes relevant when the problem is generalized to consider more than one distinguished item, or even fractional amounts thereof – so that ***a* can be any real**

**number in the unit interval.**



# Quantum Amplitude Estimation

$$|\psi\rangle = \sqrt{a} \left| \begin{array}{c} 0 \\ \text{treasure chest} \end{array} \right\rangle + \sqrt{1-a} \left| \begin{array}{c} 0 \\ \text{empty chest} \end{array} \right\rangle \quad \boxed{a = ?}$$

Classical approach: sum over samples.

$$\frac{1}{N} \sum_x f(x) = \frac{1}{N} \sum_{x \in X} 1 = \frac{1}{N} M = a$$

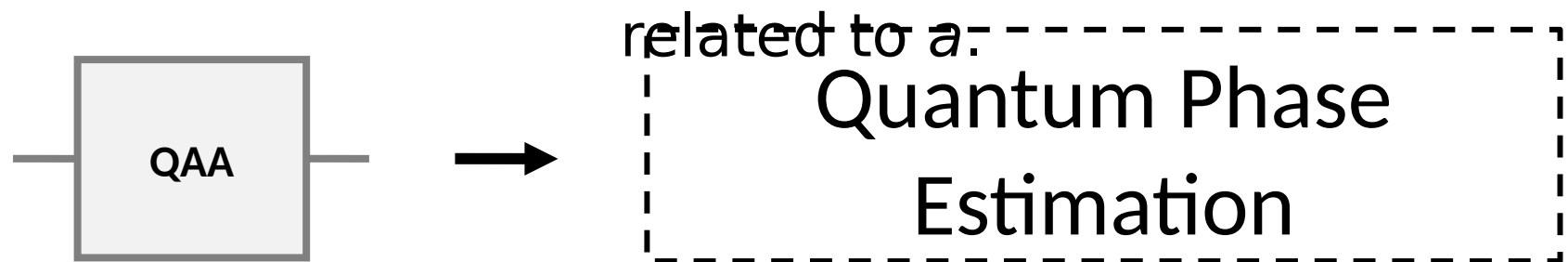
$$\epsilon \in \mathcal{O}(1/\sqrt{K})$$

# Quantum Amplitude Estimation

Quantum approach: perform **phase estimation**  
**on the amplification operator**  $\lambda_{\pm} = \exp(\pm i2\theta)$

eigenvalues

are directly



We measure one of two eigenphases: or

$$\epsilon \in O(1/K)$$

# Quantum Amplitude Estimation

$$|\psi\rangle = \sqrt{a} \left| \begin{array}{c} 0 \\ \text{treasure chest} \end{array} \right\rangle + \sqrt{1-a} \left| \begin{array}{c} 0 \\ \text{empty chest} \end{array} \right\rangle \quad \boxed{a = ?}$$

This estimation task is a fundamental routine, with a wide range of applications in **chemistry**, **machine learning** and **statistics**.

# Quantum Amplitude Estimation

In particular, it can be used to **speed-up Monte**

$$\mathbb{E}_{p(x)}[f(x)] = \int_{\Omega} f(x)p(x)dx$$

$$\mathbb{E}_{p(x)}[f(x)] \approx \frac{1}{N} \sum_{i=0}^{N-1} \frac{f(x_i) \cdot p(x_i)}{\pi(x_i)} \quad \{x_i\}_{i=0}^{N-1} \sim \pi(\cdot)$$

$$i \sum_x p(x) f(x) \quad \text{for a uniform PMF}$$

# Quantum Amplitude Estimation

Define **distribution loading** and **function encoding**

$$P|0\rangle^{\otimes n} = \sum_{x=0}^{2^n-1} \sqrt{p(x)} |x\rangle \quad R|x\rangle|0\rangle = |x\rangle \left( \sqrt{f(x)} |1\rangle + \sqrt{1-f(x)} |0\rangle \right)$$

$$R(P \otimes I_1)|0\rangle^{\otimes n} = \sum_{x=0}^{2^n-1} \sqrt{p(x)} |x\rangle \left( \sqrt{f(x)} |1\rangle + \sqrt{1-f(x)} |0\rangle \right)$$

Rewrite in terms of orthonormal subspaces and normalize to get:

$$|\psi\rangle = |\psi_1\rangle + |\psi_0\rangle = \sqrt{a} |\tilde{\psi}_1\rangle + \sqrt{1-a} |\tilde{\psi}_0\rangle$$

$$a \equiv \sum_x p(x) f(x) = \mathbb{E}_{p(x)}[f(x)]$$

We have an amplitude estimation problem!



# Quantum Amplitude Estimation

Define **distribution loading** and **function encoding**

$$P|0\rangle^{\otimes n} = \sum_{x=0}^{2^n-1} \sqrt{p(x)} |x\rangle \quad R|x\rangle|0\rangle = |x\rangle \left( \sqrt{f(x)} |1\rangle + \sqrt{1-f(x)} |0\rangle \right)$$

$$R(P \otimes I_1)|0\rangle^{\otimes n} = \sum_{x=0}^{2^n-1} \sqrt{p(x)} |x\rangle \left( \sqrt{f(x)} |1\rangle + \sqrt{1-f(x)} |0\rangle \right)$$

Rewrite in terms of orthonormal subspaces:

$$|\psi_1\rangle = \sum_{x=0}^{2^n-1} \sqrt{p(x)} \sqrt{f(x)} |x\rangle |1\rangle \quad |\psi_0\rangle = \sum_{x=0}^{2^n-1} \sqrt{p(x)} \sqrt{1-f(x)} |x\rangle |0\rangle$$

Normalize to get:  $|\psi\rangle = |\psi_1\rangle + |\psi_0\rangle = \sqrt{a} |\tilde{\psi}_1\rangle + \sqrt{1-a} |\tilde{\psi}_0\rangle$

We have an amplitude estimation problem!

# Quantum Amplitude Estimation

$$a \equiv \sum_x p(x) f(x) \approx \mathbb{E}_{p(x)} [f(x)]$$

Our **initialization** and **oracle** operators are now:

$$A \equiv R(P \otimes I_1)$$

$$\hat{U}_j = (I_n \otimes Z)$$

*Example:* integrate  $\frac{1}{b_{\max}} \int_0^{b_{\max}} \sin^2(x) dx$  using a uniform importance distribution.

using a uniform

$$R |x\rangle |0\rangle = |x\rangle \left( \underbrace{\sin\left(\frac{(x + \frac{1}{2})b_{\max}}{2^n}\right)}_{\frac{b_{\max}}{2^n} + \sum_{k=1}^n \frac{b_{\max} \cdot x_k}{2^{n-k}} / 2} |1\rangle + \cos\left(\frac{(x + \frac{1}{2})b_{\max}}{2^n}\right) |0\rangle \right) \rightarrow R = \prod_{k=1}^n C^{(k)} R_y\left(\frac{b_{\max} \cdot x_k}{2^{n-k}}\right) \left( I_n \otimes R_y\left(\frac{b_{\max}}{2^n}\right) \right)$$

$$P = H^{\otimes n}$$

# Quantum Amplitude Estimation

we can calculate the measurement probability for any bit string

$$P(\text{measuring } x \mid QAE(\theta)) = \frac{P(\text{measuring } x \mid QPE(\theta/\pi)) + P(\text{measuring } x \mid QPE(1 - \theta/\pi))}{2}.$$

*with*

$$P(\text{measuring } x \mid QPE(\phi)) = \frac{\sin^2(K\Delta\pi)}{K^2 \sin^2(\Delta\pi)},$$

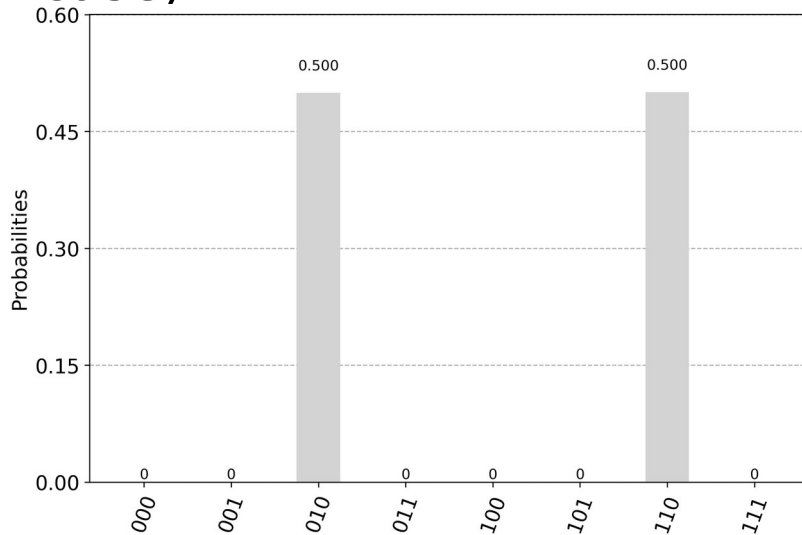
where  $\Delta$  is a circular distance, and also the error in the estimate produced by  $x$ :

$$\Delta = \left| \phi - \frac{x}{K} \right| \pmod{1}.$$

# Quantum Amplitude Estimation

When  $x$  is an integer, we retrieve the exact solution by measuring.

Example bar plot (deterministic case):



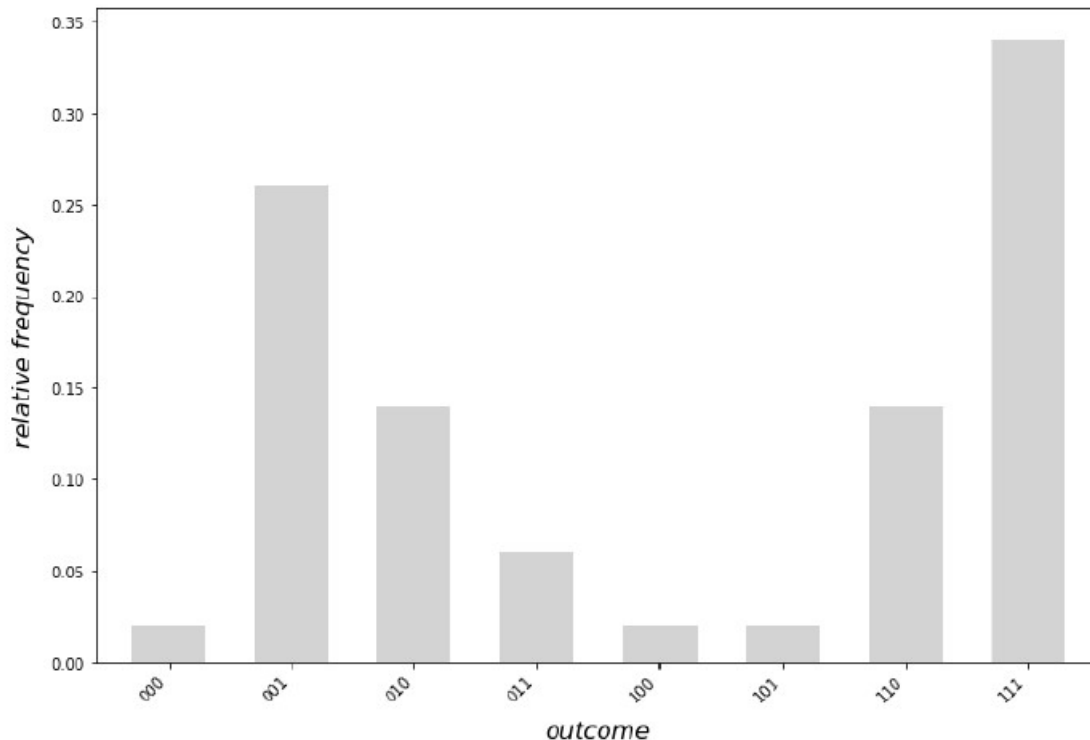
- 50 measurements

$$x=0b010=2 \rightarrow \theta = \frac{x\pi}{K} = \frac{2\pi}{8}$$

$$x=0b110=6 \rightarrow \theta = \frac{x\pi}{K} = \frac{6\pi}{8}$$

# Quantum Amplitude Estimation

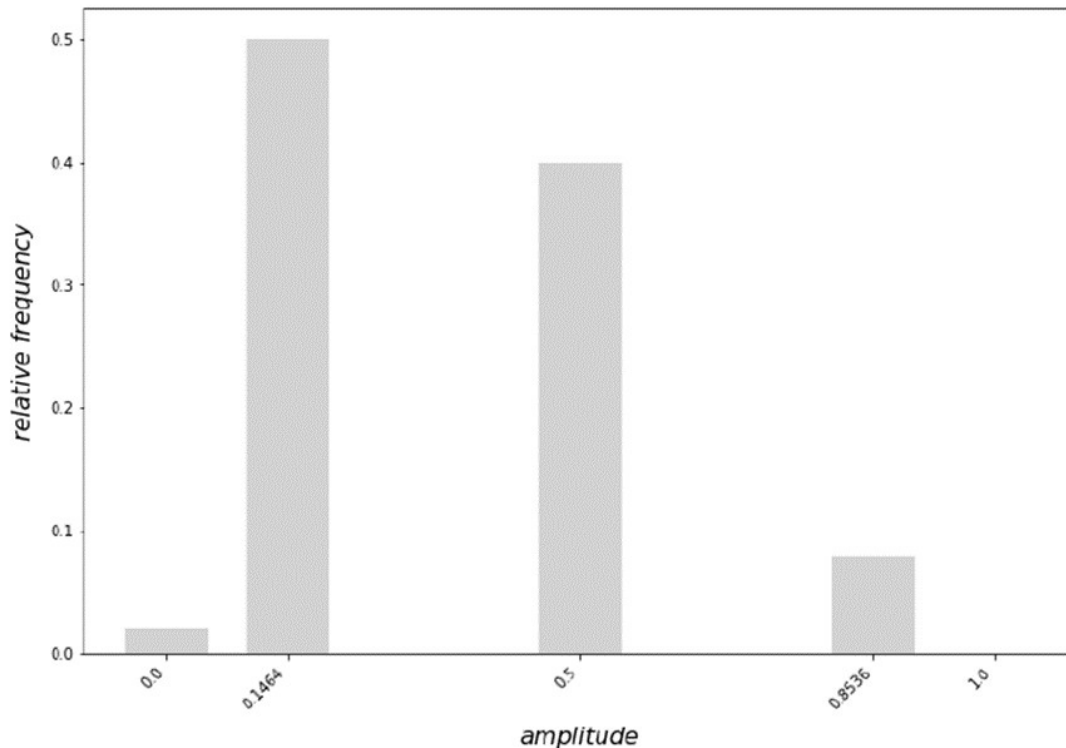
When  $\theta$  is not an integer, the outcome distribution is not so neat.  
Example bar plot:



50 measurements

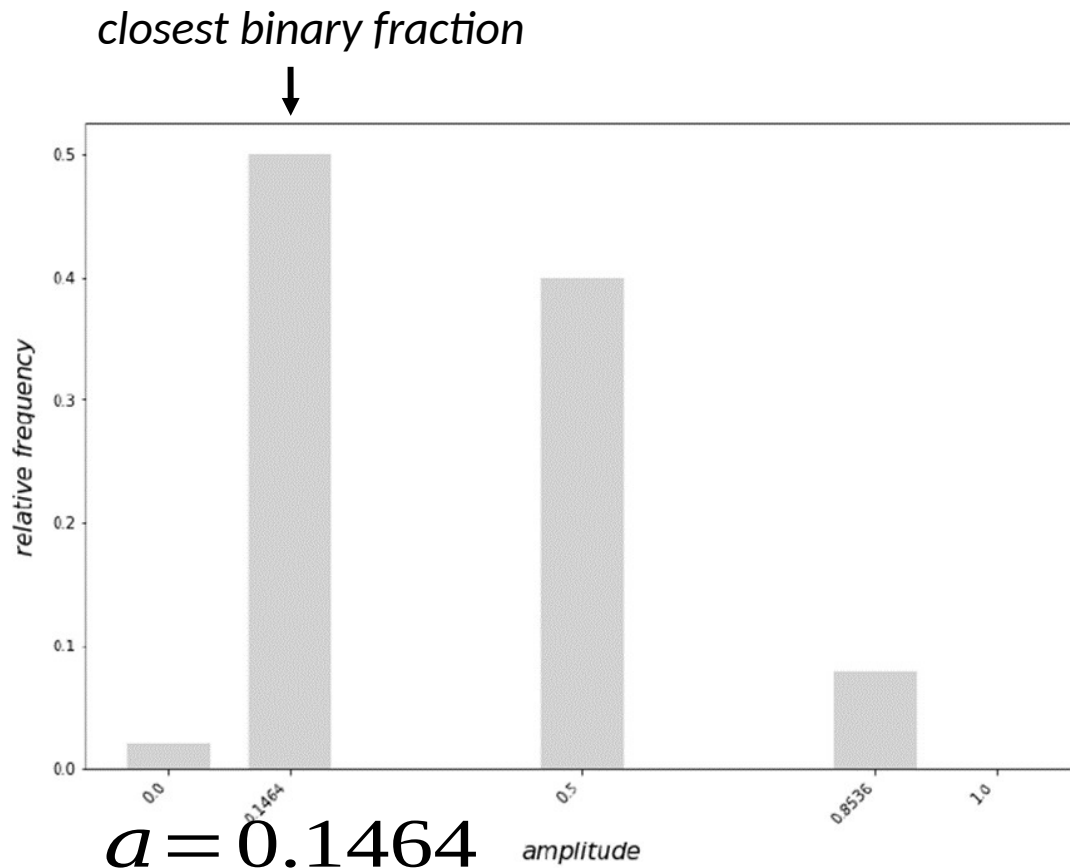
# Quantum Amplitude Estimation

We can still translate it into the amplitude domain:



- 50 measurements

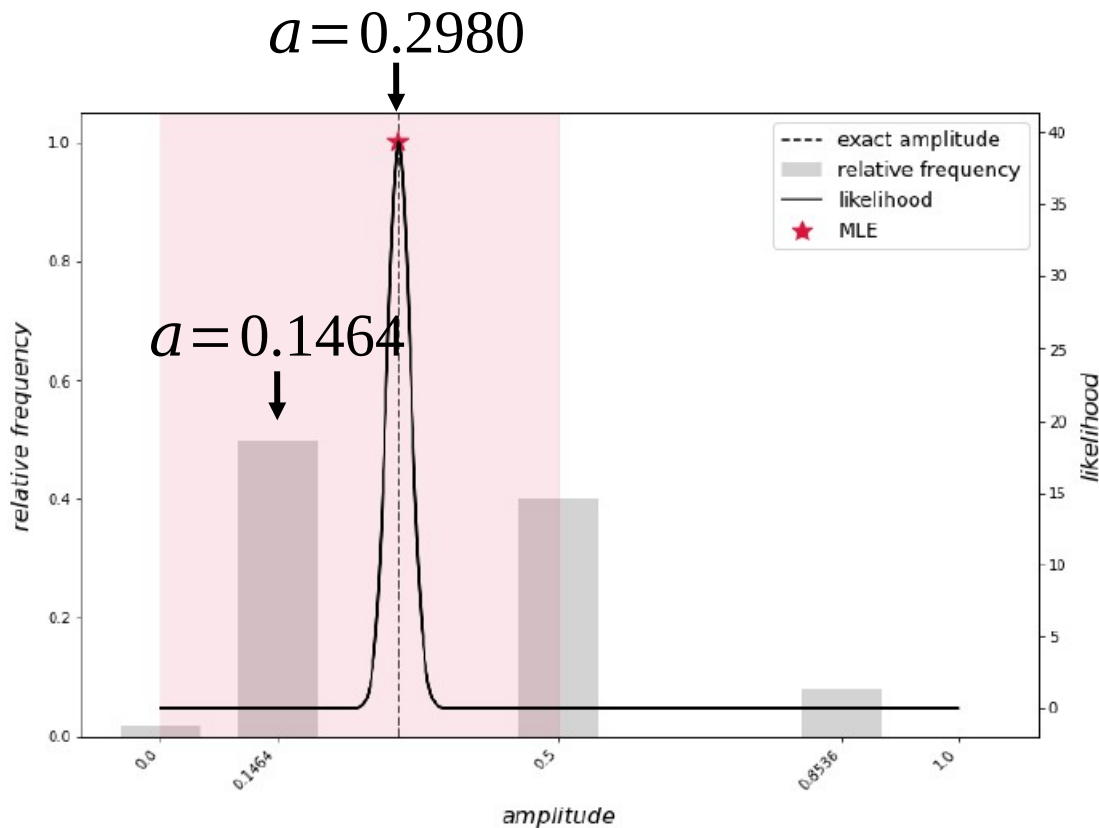
# Quantum Amplitude Estimation



- 50 measurements

# Quantum Amplitude Estimation

Rather than sticking to a grid, we can use our knowledge of the outcome distribution to sweep over a continuum of values for  $a$ .

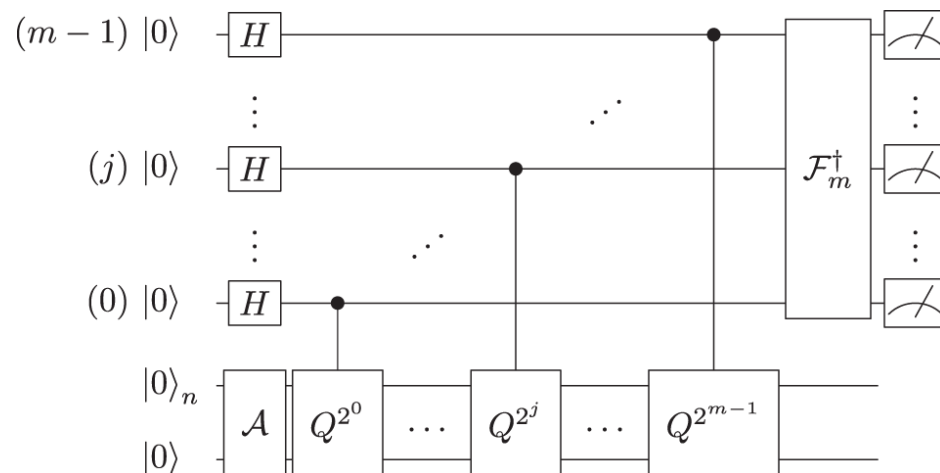


- 50 measurements



# Quantum Amplitude Estimation

Today's quantum devices are unable to gainfully realize **QPE**.

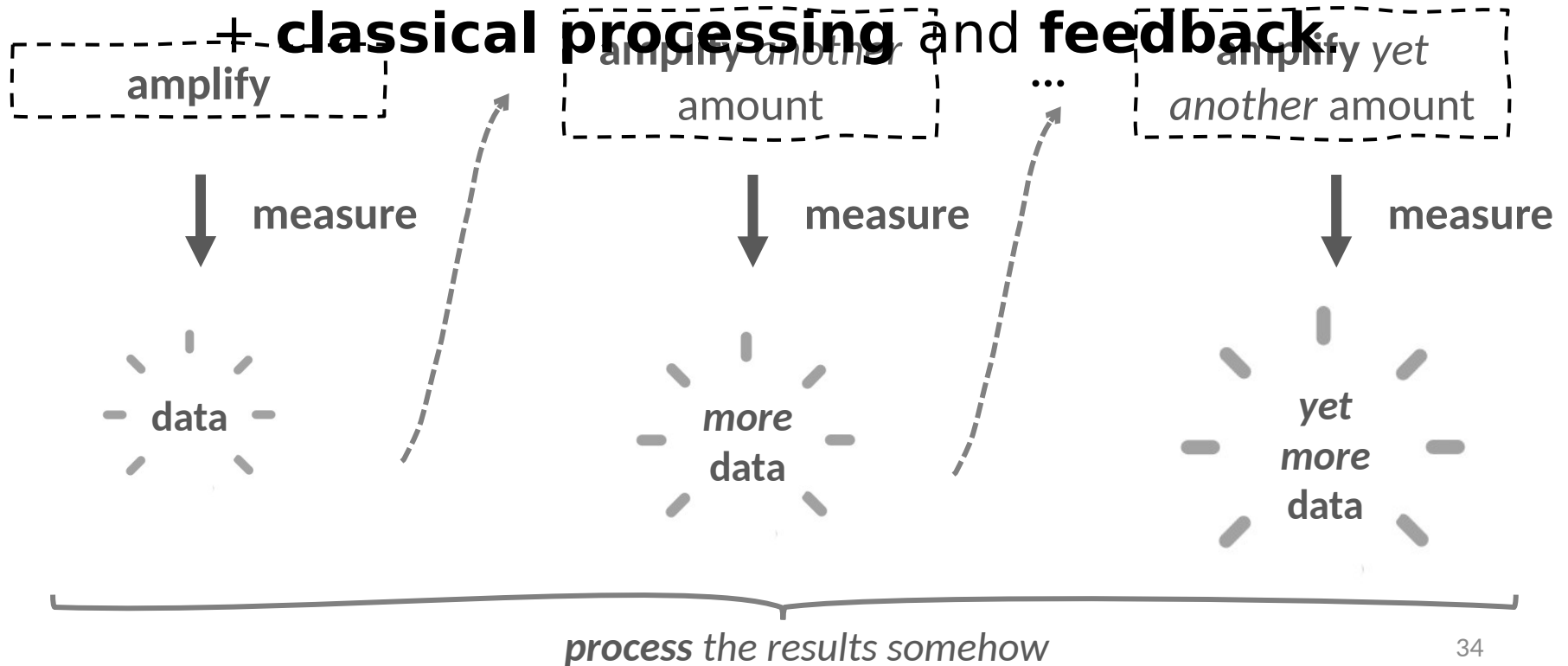


*in: Grinko, D., Gacon, J., Zoufal, C. et al. Iterative quantum amplitude estimation. npj Quantum Inf 7, 52 (2021).*

As such, **alternative strategies** have been proposed to achieve **quantum-enhanced precision** in more **hardware-friendly** ways.

# Quantum Amplitude Estimation

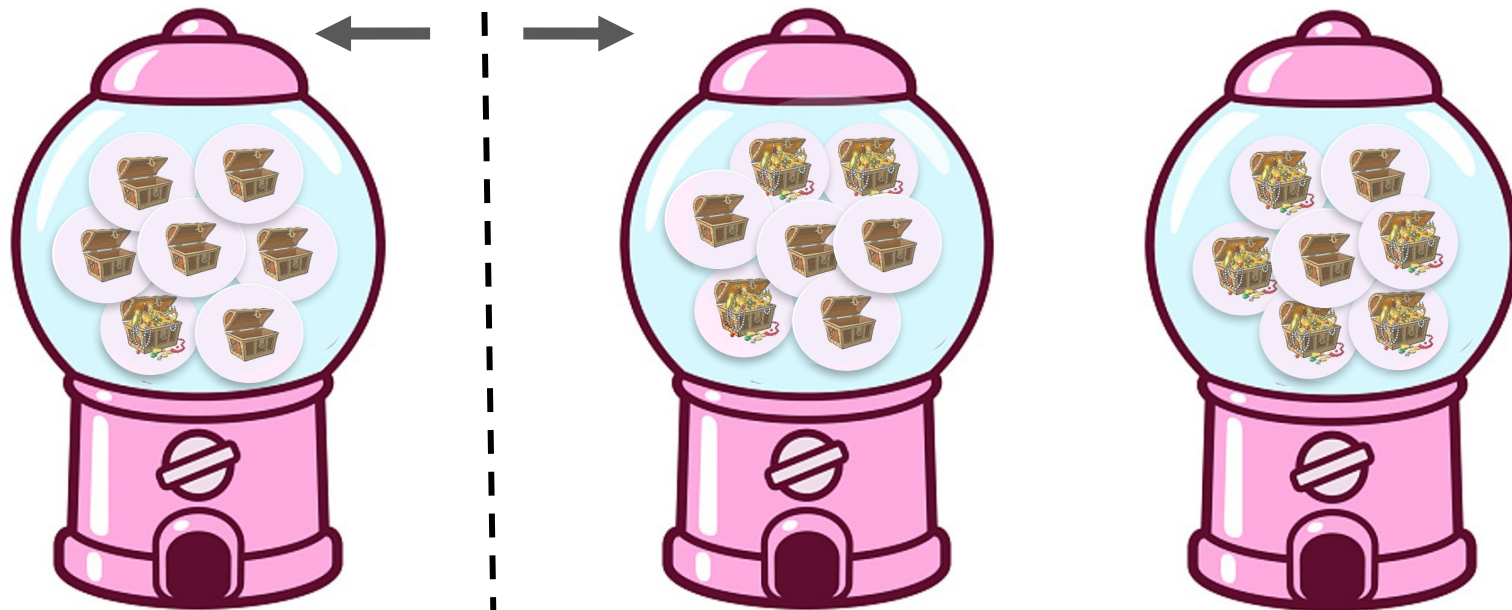
These approaches are often **hybrid** and **adaptive**, relying on **simpler quantum circuits**



# Quantum Amplitude Estimation

The overarching idea is to use **quantum amplitude amplification** to make **more**

**informative** measurements.



# Quantum Amplitude Estimation

$$\text{Binomial}(p(\theta), N_{\text{shots}})$$

**Classically**, we can  
sample from:

**Quantum-enhanced**  
measurements allow:

*for any odd integer*

# Quantum Amplitude Estimation

With ,

Preparing entails queries to  $A$  (forwards or backwards):

- for the initialization of
- for the applications of

It also requires queries to the oracle, one for each application of .

We can sample according to the probability

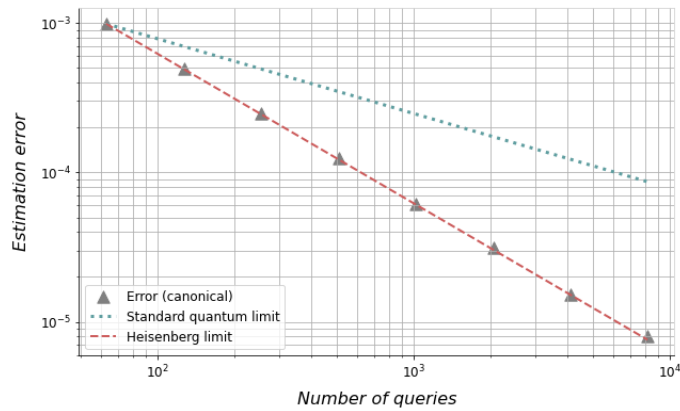
**using a number of queries is in .**

*How* to use these measurements is an open question.  
Each answer brings a different **quantum advantage**,  
**classical overhead** and **noise resilience**.

***Overview and numerical  
analysis of QAE  
algorithms***

# Textbook QAE

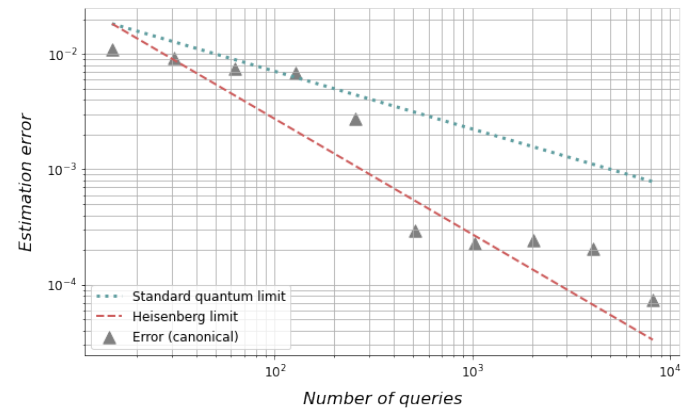
Scaling of the estimation error in  $a$  with the number of queries to  $A$



$$a = 0.5$$

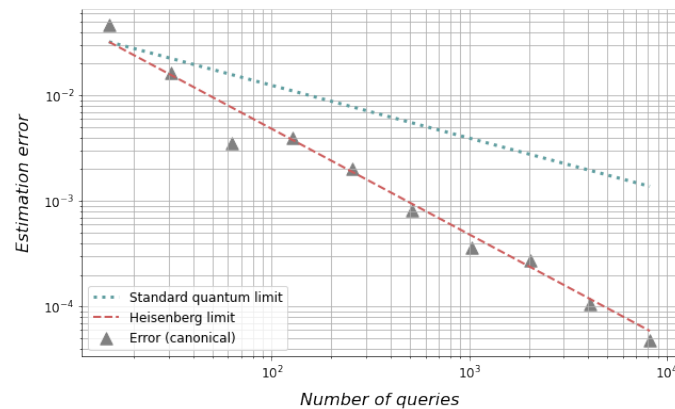
$$\epsilon^{(HL)} \propto \frac{1}{N}$$

Scaling of the estimation error in  $a$  with the number of queries to  $A$



$$a = \frac{\pi}{10}$$

Scaling of the estimation error in  $a$  with the number of queries to  $A$



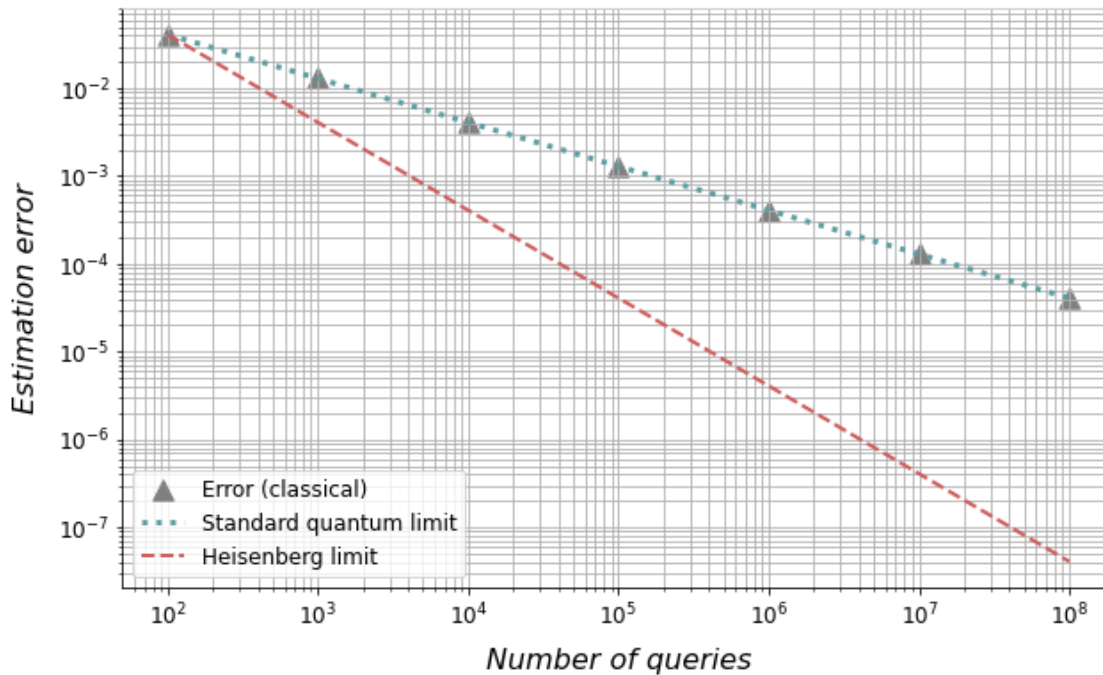
$$a \text{ Unif } i$$



# Classical Amplitude Estimation

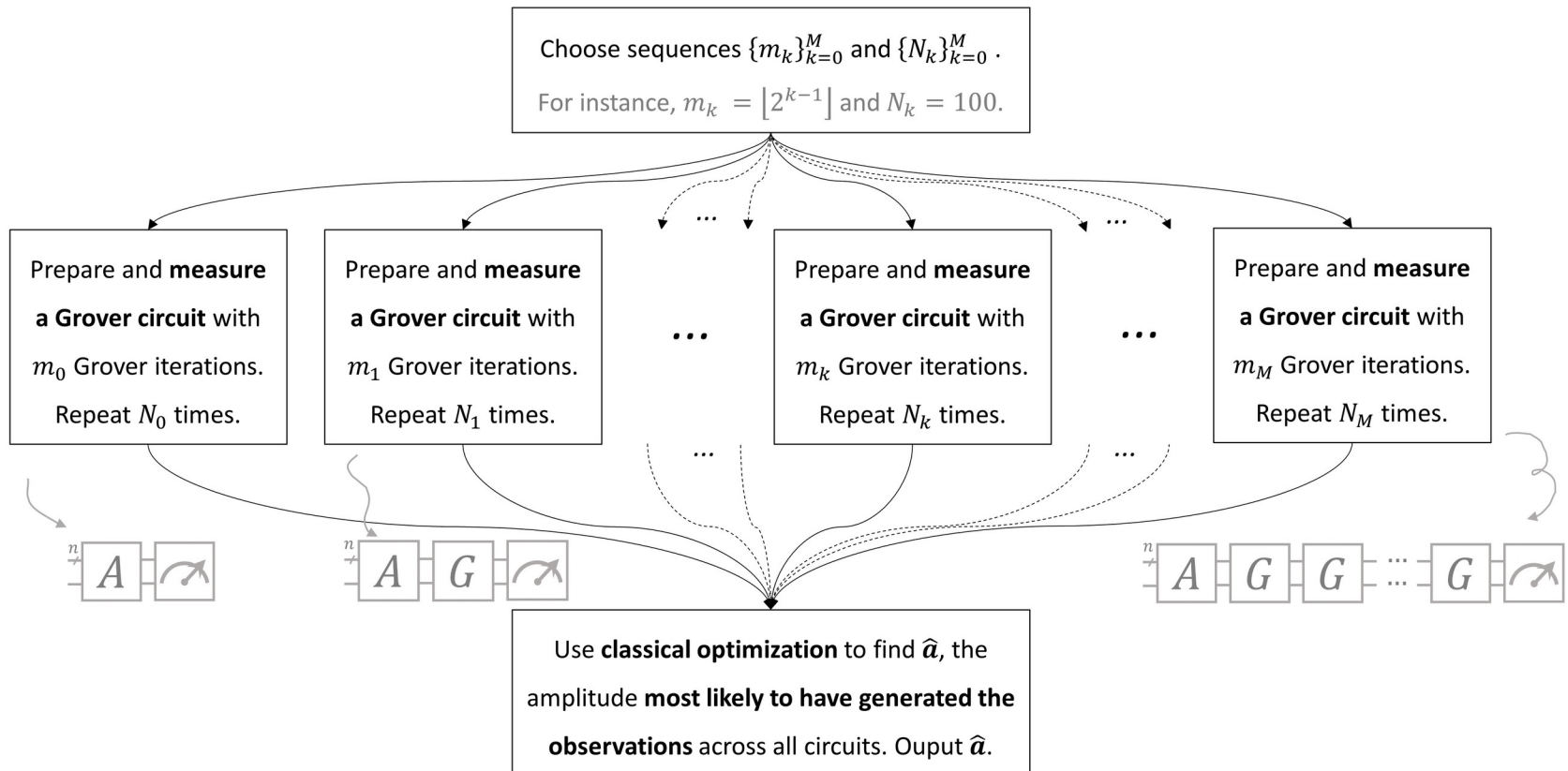
Compare with the classical case:

Scaling of the estimation error in a with the number of queries to A



$$\epsilon^{(SQL)} \propto \frac{1}{\sqrt{N}}$$

# Maximum Likelihood QAE

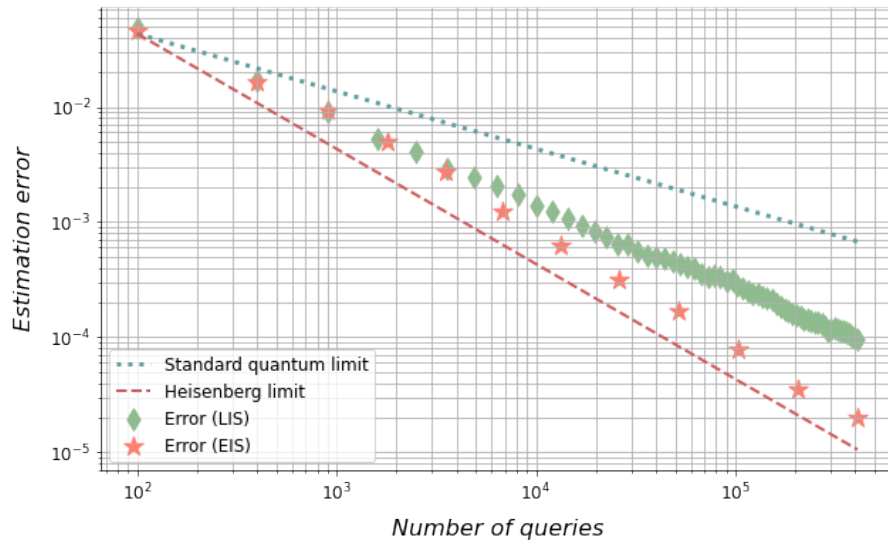


# Maximum Likelihood QAE

- Original proposal (Suzuki et al): ;
- Spin offs:
  - Variational QAE
  - Low depth QAE  $N \in \mathcal{O}(1/\epsilon^{1+\beta})$  and  $D \in \mathcal{O}(1/\epsilon^{1-\beta})$ 
    - Power Law: , a controllable parameter
    - QoPrime

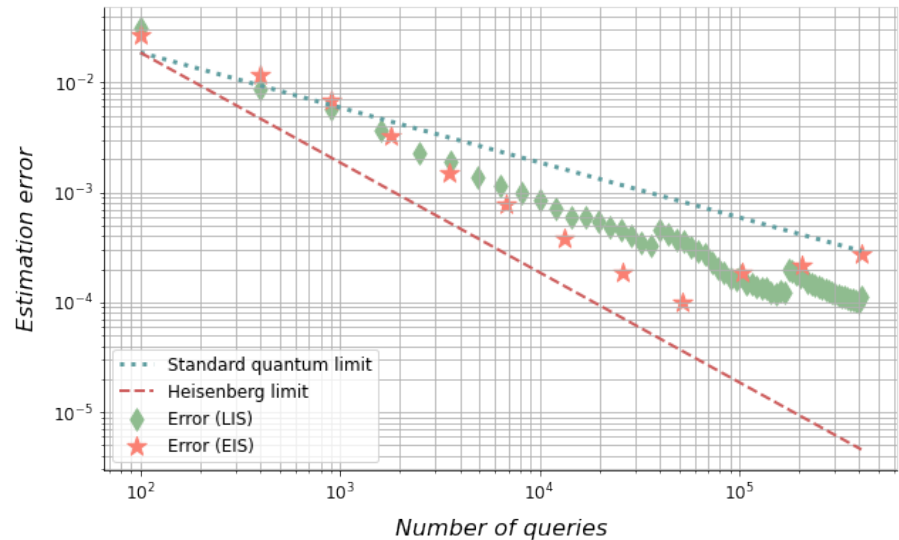
# Maximum Likelihood QAE

Scaling of the estimation error in a with the number of queries to A



$$a = 0.5$$

Scaling of the estimation error in a with the number of queries to A



# Quantum Amplitude Estimation

## A note on performance benchmarking for adaptive algorithms

- Sample an  $x$  (representing  $N$ ) coordinate uniformly at random on a log scale.

$$x \sim \text{unif}_{\log}([x_{\min}, x_{\max}])$$

- Sample auxiliary variables  $z$  depending on  $x$ :

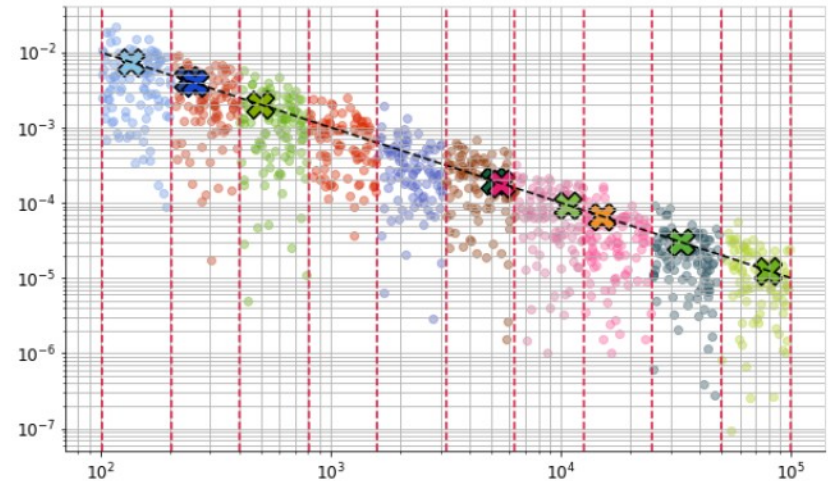
$$z \sim \mathcal{N}(\mu, \sigma(x)),$$

where

$$\sigma(x) \propto 1/x.$$

The choice of the mean  $\mu$  is arbitrary, whereas the constant of proportionality is determined by fixing a point.

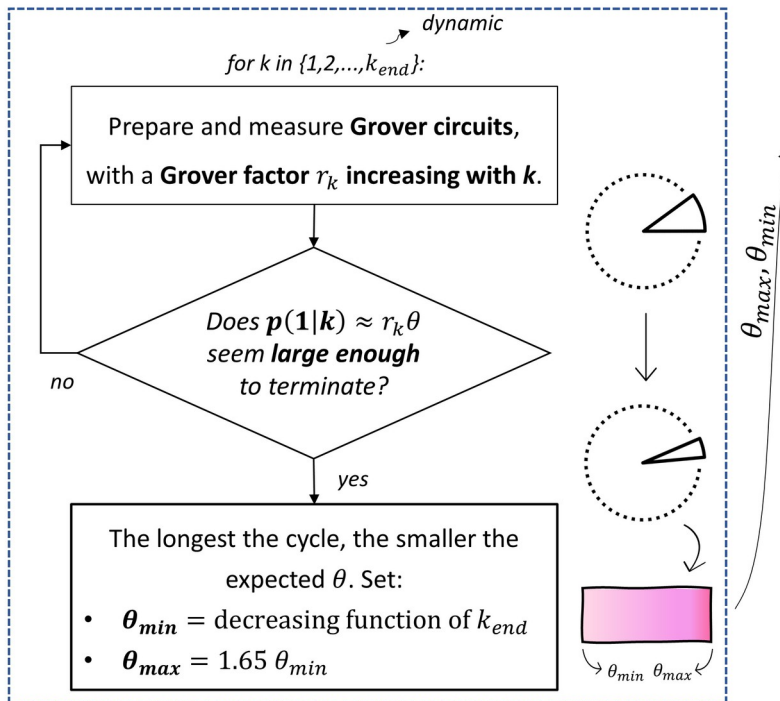
- Calculate  $y$  (representing  $\epsilon^2$ ) as  $y = (x - \mu)^2$ .



# QAE, simplified

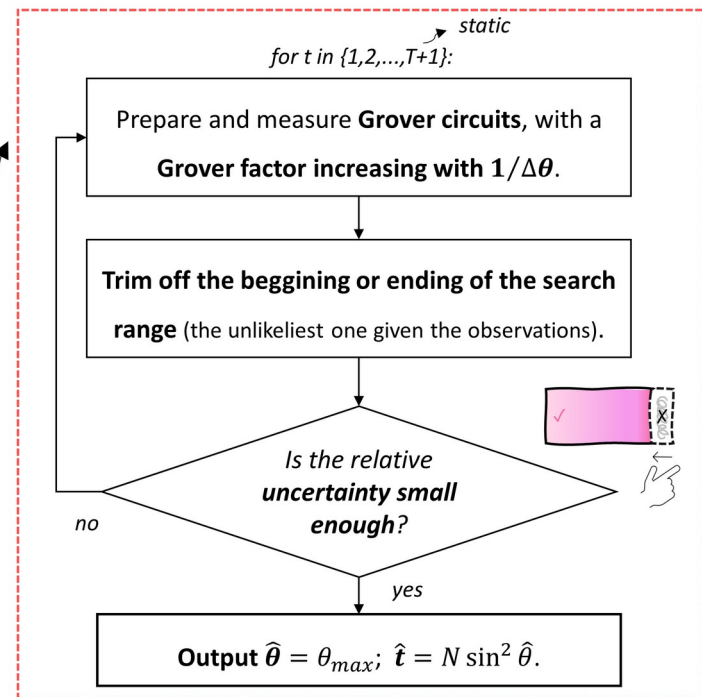
## Pre-processing phase

Reduce the search range to  $[\theta_{min}, \theta_{max}]$ , with  $\frac{\theta_{max}}{\theta_{min}} = 1.65$ .



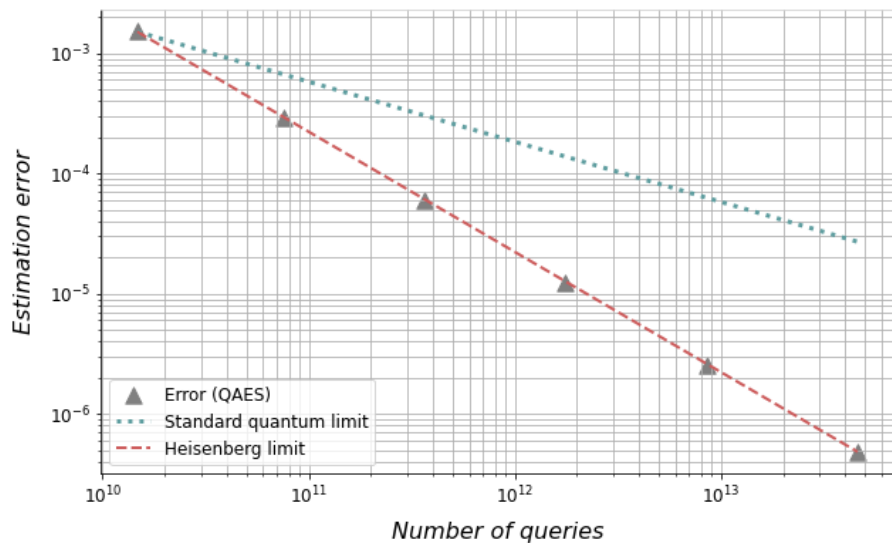
## Exponential refinement phase

Reduce the relative uncertainty by 10% per iteration.

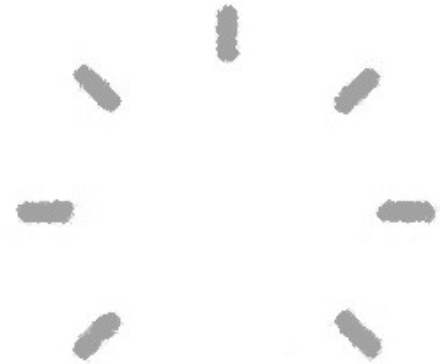


# QAE, simplified

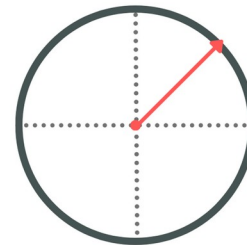
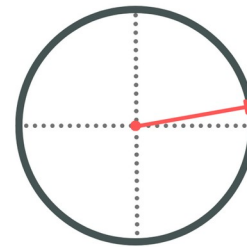
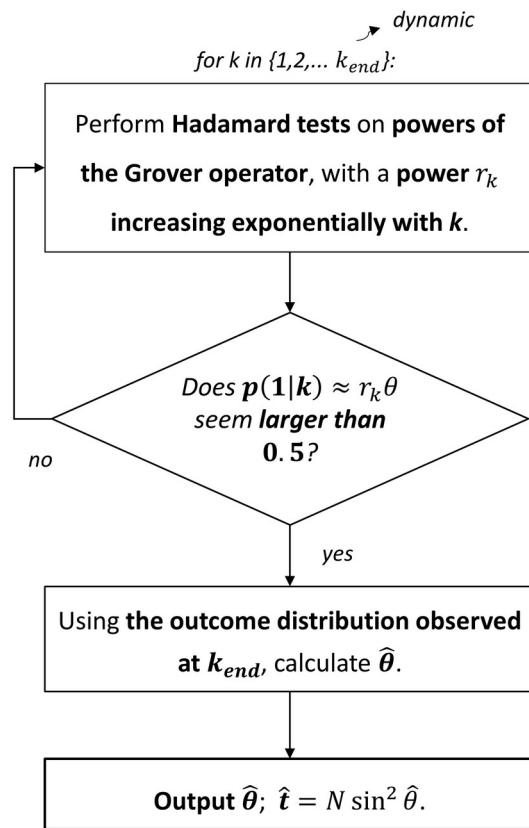
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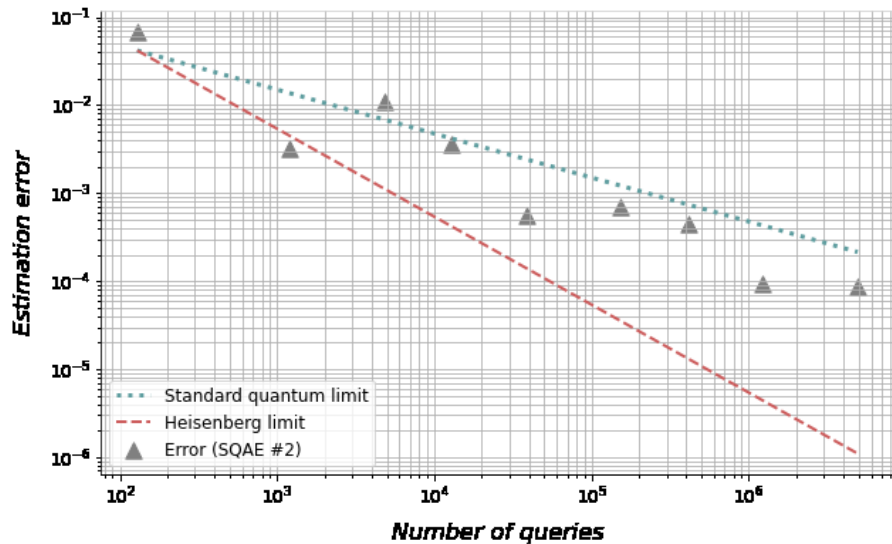
# Simpler QAE





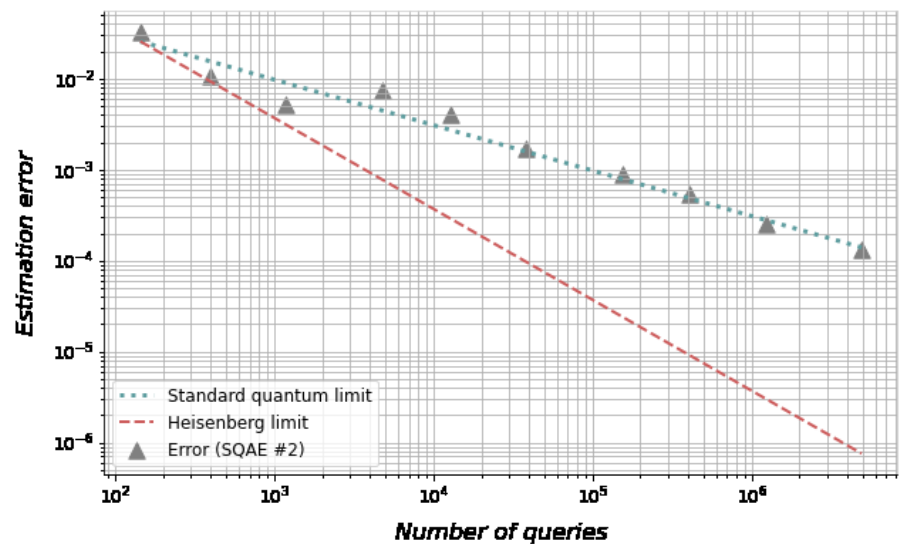
# Simpler QAE

Scaling of the estimation error in  $a$  with the number of queries to  $A$



$$a = 0.5$$

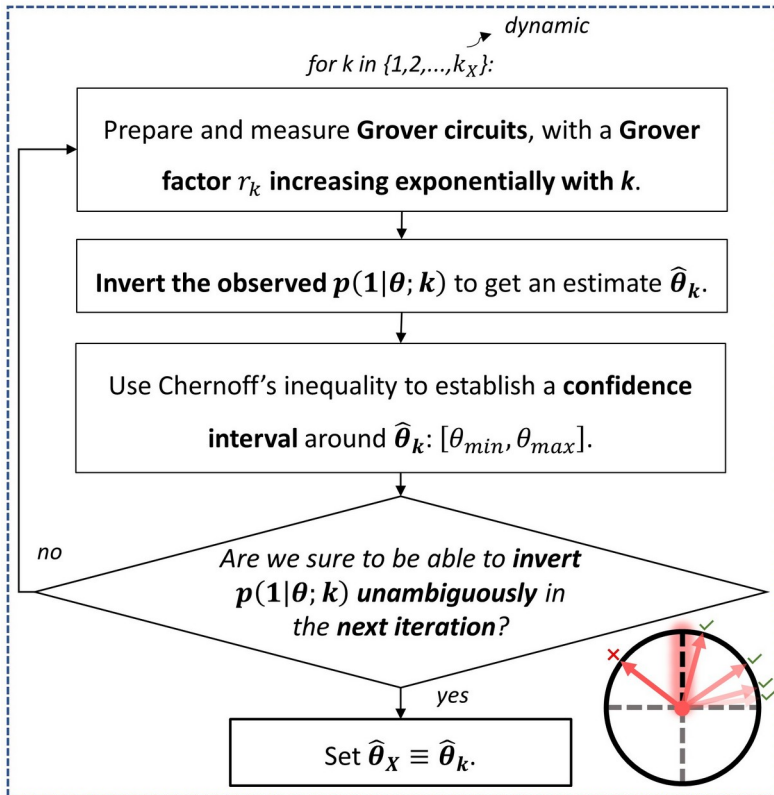
Scaling of the estimation error in  $a$  with the number of queries to  $A$



# Faster QAE

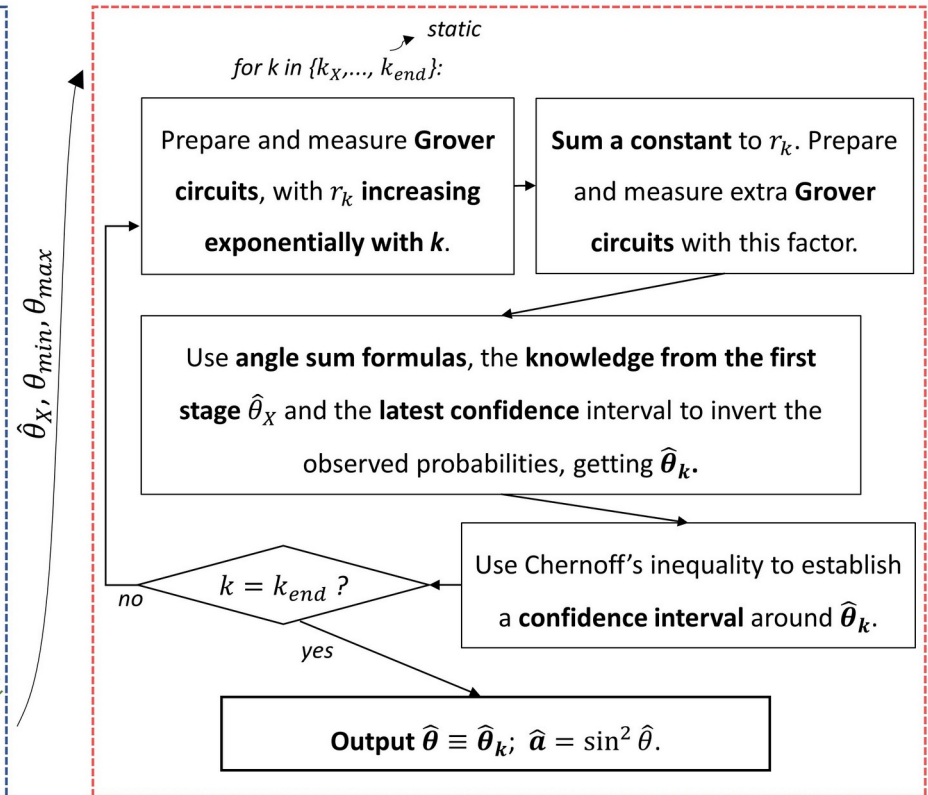
## First stage

Learn by **inverting** the results of **Grover measurements** with **increasing amplification**, until they **become ambiguous**.



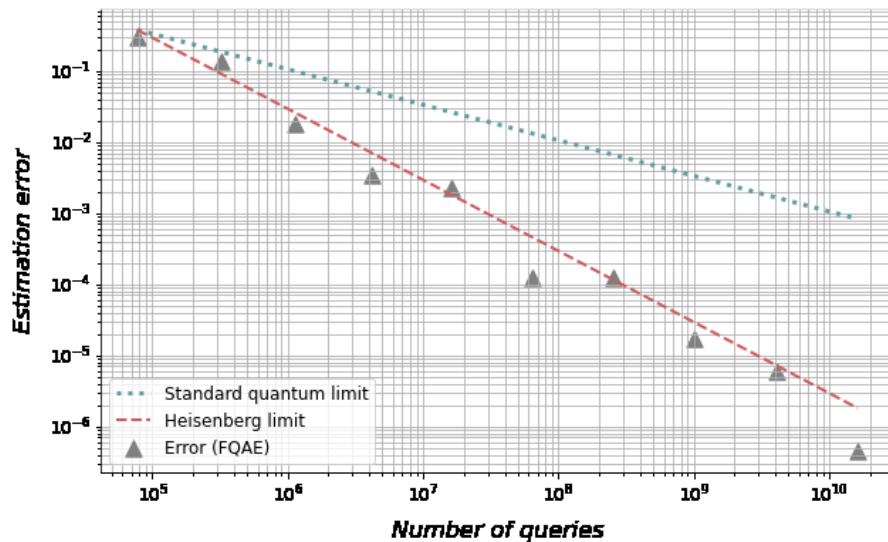
## Second stage

Continue to learn by **complementing the former Grover measurements** with additional ones for **disambiguation**.



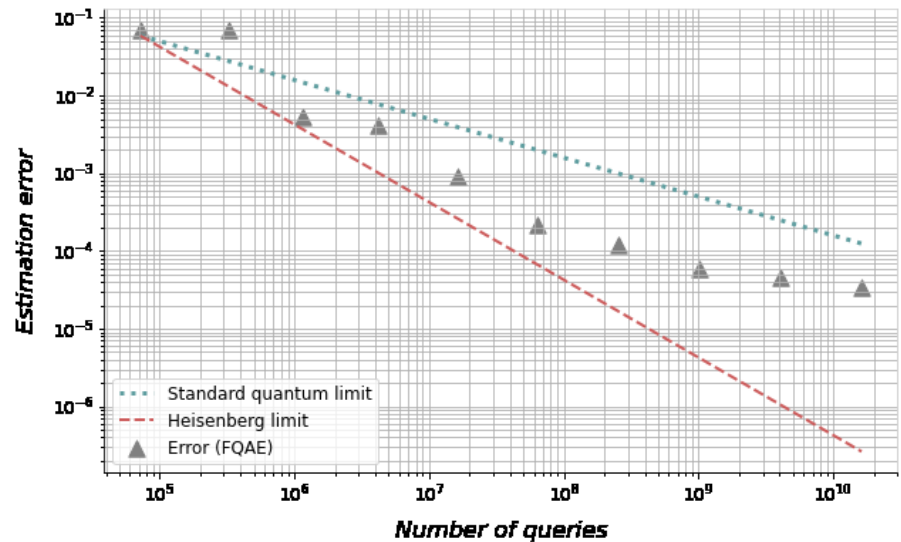
# Faster QAE

Scaling of the estimation error in  $a$  with the number of queries to  $A$

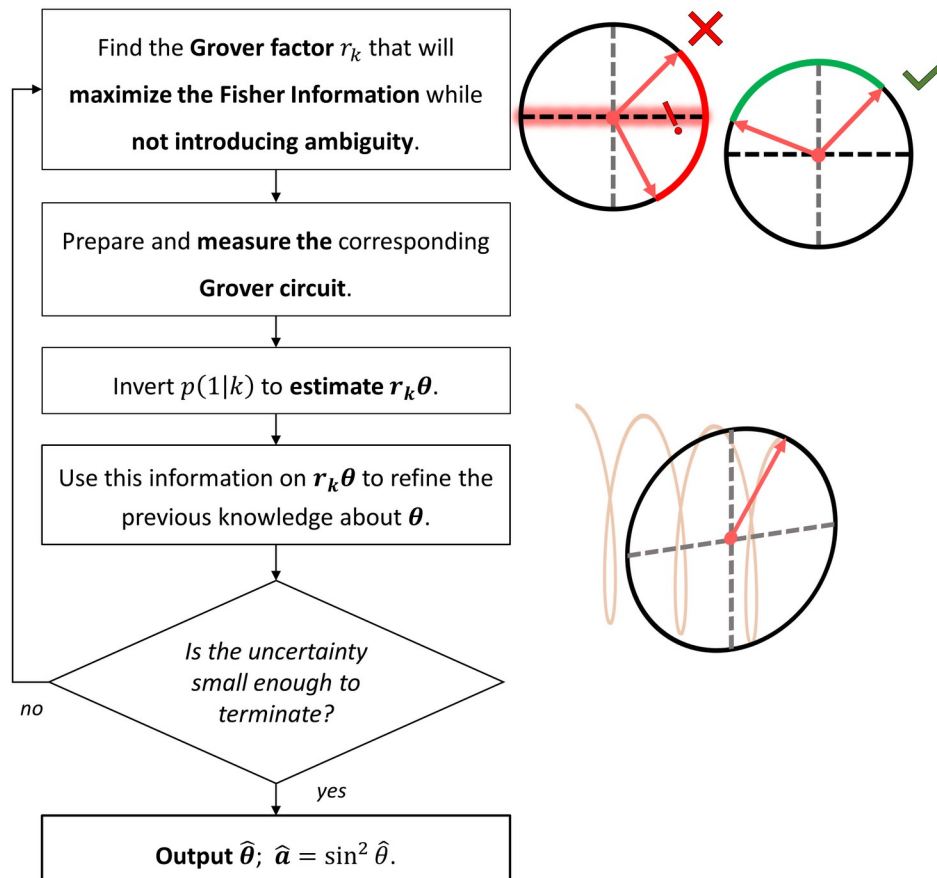


$$a = 0.5$$

Scaling of the estimation error in  $a$  with the number of queries to  $A$

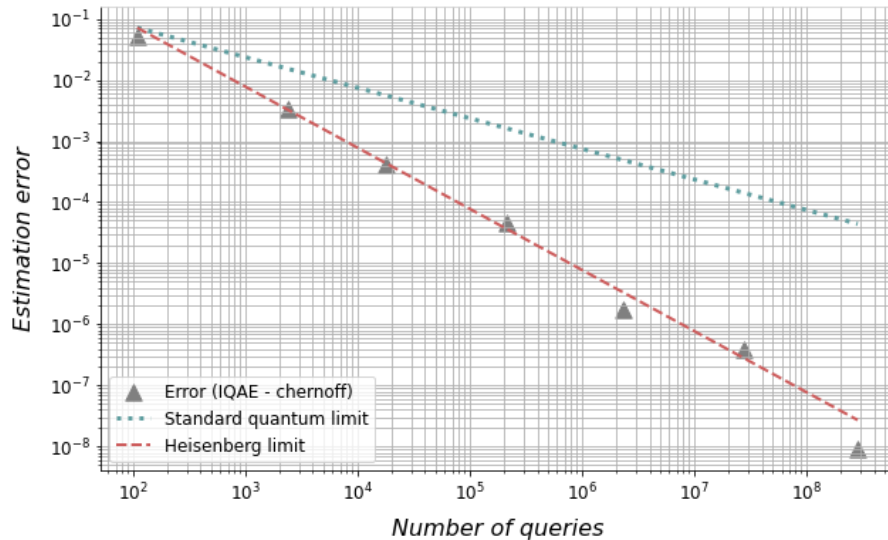


# Iterative QAE



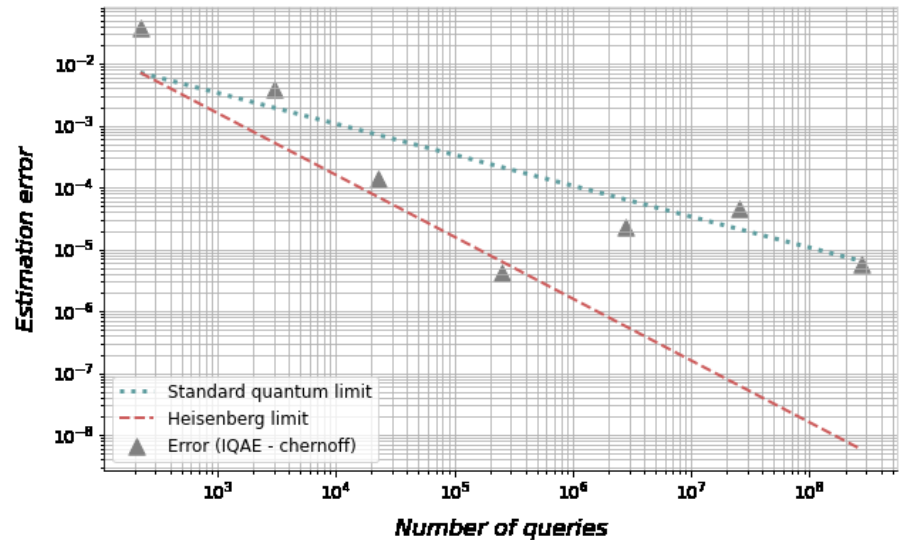
# Iterative QAE

Scaling of the estimation error in  $a$  with the number of queries to  $A$



$$a = 0.5$$

Scaling of the estimation error in  $a$  with the number of queries to  $A$

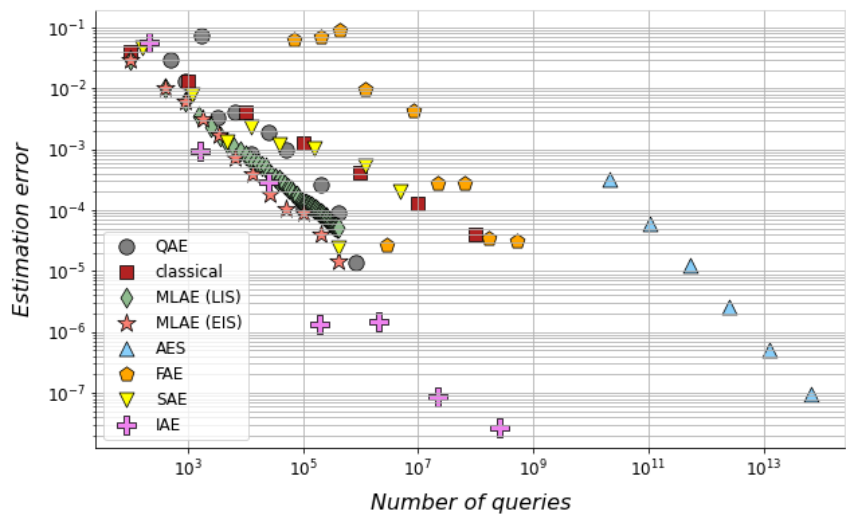


# Summary

	Key idea	Parallelizable	Circuits	Main strength(s)	Main weakness(es)	Complexity
<b>QAE</b> [4]	QPE on Grover operator	no	QPE	optimal complexity	very deep circuit; noise oblivious	$N_q \in O(\frac{1}{\epsilon} \log \frac{1}{\alpha} \cdot a^{-1})$
<b>MLAE</b> [56]	heuristic measurements $\rightarrow$ statistical estimation	fully	QAA	simplicity; solid numerical performance	no formal guarantees; unchecked circuit growth	best case   observed: $N_q^{(US)} \sim \epsilon^{-0.75} \mid \epsilon^{-0.74}$ $N_q^{(ES)} \sim \epsilon^{-1} \mid \epsilon^{-0.88}$
<b>V-MLAE</b> [47]	MLAE with variational approximation	fully	QAA	solid numerical performance; limits circuit depth	no formal guarantees; cost overhead	similar to MLAE except for EIS-observed (untested)
<b>AES</b> [1]	rough localization $\rightarrow$ exponential refinement	partly	QAA	optimal complexity	large cost offset; noise oblivious	$N_q \in O(\frac{1}{\epsilon} \log \frac{1}{\alpha} \cdot a^{-1})$
<b>SAE</b> [61]	amplify $\rightarrow$ invert probability	partly	Hadamard tests	—	inconclusive demonstrations	$N_q \in O(a^{-1})$
<b>IAE</b> [22]	watchful choice of Fisher information	partly	QAA	nearly optimal complexity	unchecked circuit growth; noise oblivious	$N_q \in O(\frac{1}{\epsilon} \log(\frac{1}{\alpha} \log \frac{1}{\epsilon}))$
<b>M-IAE</b> [17]	IAE but distribute shots more favorably	partly	QAA	optimal complexity	same as IAE	$N_q \in O(\frac{1}{\epsilon} \log \frac{1}{\alpha})$
<b>FAE</b> [42]	complementary measurements $\rightarrow$ invert probability	partly	QAA	nearly optimal complexity	unchecked circuit growth; noise oblivious	$N_q \in O(\frac{1}{\epsilon} \log(\frac{1}{\alpha} \log \frac{1}{\epsilon}))$

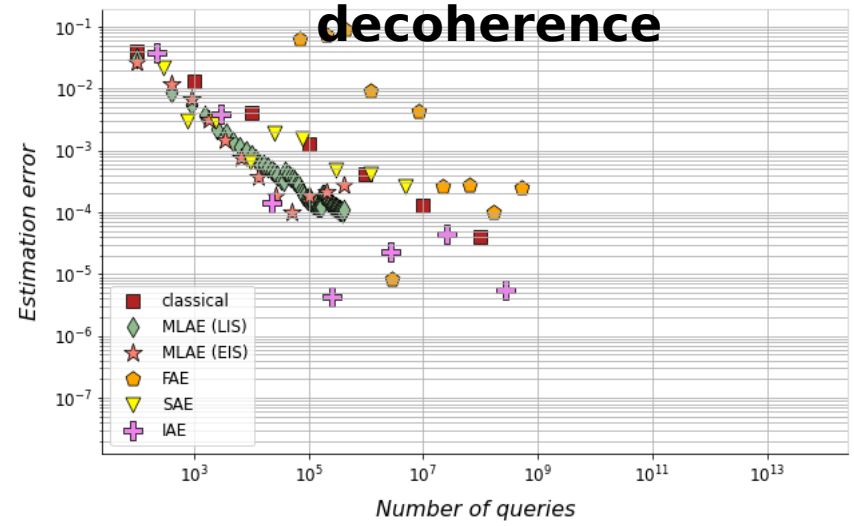
# Summary

## ideal performance



$a = 0.1$

## performance in the presence of decoherence



# Bayesian amplitude estimation

## Bayesian inference

offers a natural paradigm for solving this problem.

$$P(\theta | D) = \frac{L(\theta | D; E)P(\theta)}{P(D; E)}$$

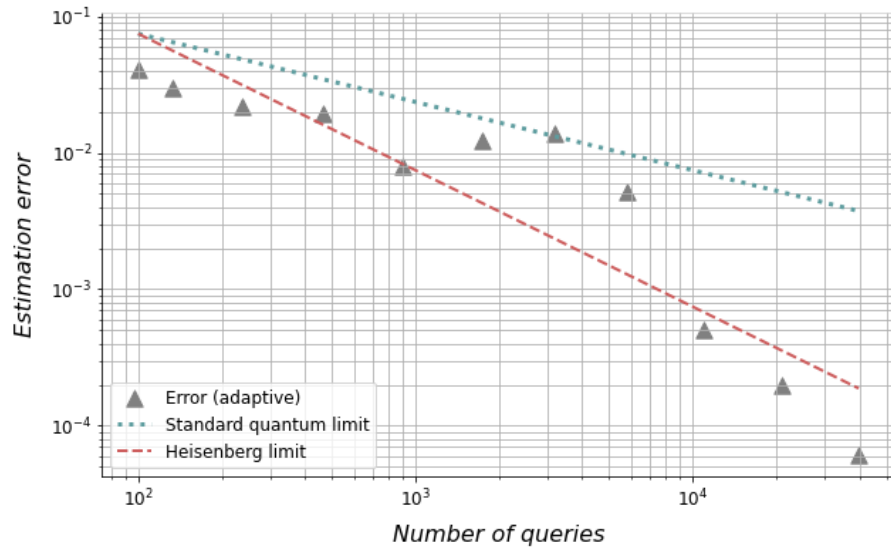
Unlike less general protocols, it inherently

- The possibility of incorporating **noise** offers:  
**models;**
- **Flexibility** to tailor the protocol;
- Various ways of negotiating the **trade-offs**



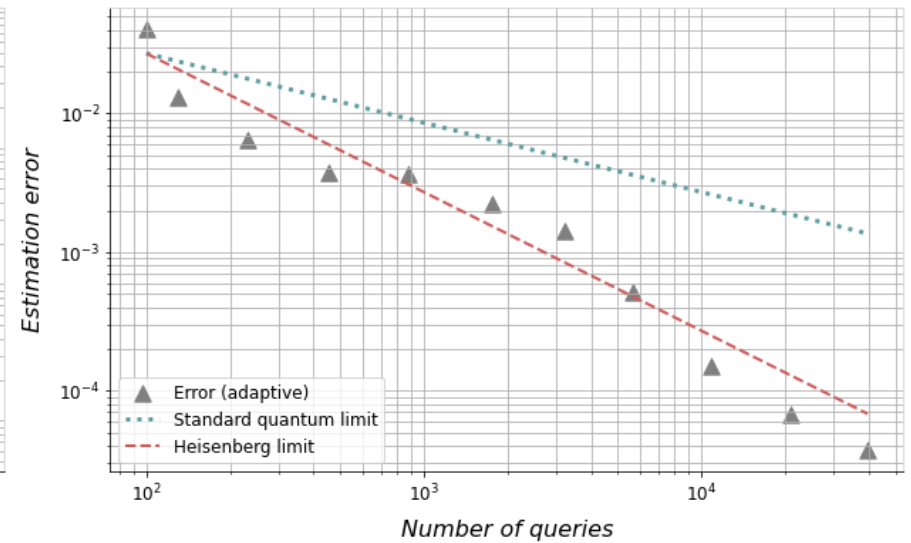
# Bayesian amplitude estimation

Scaling of the estimation error in  $a$  with the number of queries to  $A$



mean  
results

Scaling of the estimation error in  $a$  with the number of queries to  $A$



median  
results

$$a=1$$

# Bayesian amplitude estimation

We can make the most of the open-endedness of Bayesian inference by:

- Properly handling **noise**;
- Using the **flexibility** to our benefit;
- Seeking the best **cost-to-benefit** ratio (across several types of *benefit* – quantum enhancement, noise resilience – and *cost* – classical processing, online processing, optimization, discretization, quantum resources, quantum depth, overheads).

# Robust amplitude estimation

- Bayesian inference with engineered likelihood functions
- Greedy variance reduction (adaptively minimize proxies thereof)
- Work under a Gaussian assumption (making the representation analytically tractable)
- Simple noise model

# Conclusion

- All **QAE algorithms suffer with** the presence of **noise**, but they **are affected differently**: strategies with better idealized behavior may actually do worse once noise is introduced.
- We can expect that different **algorithms respond differently to different types of noise**, both qualitatively and quantitatively.
- Often, **algorithm design focuses on “hardware-friendly” circuits**. To make the most out of the available resources, the **processing should be “hardware-friendly”** as well.

**Thank you for your  
attention!**

*Questions?*