

# Multi-Party Computation: Providing a Post-Quantum Solution (part 1)

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- 1. Basic notions
- 2. Symmetric & Assymetric Encryption
- 3. Impact of Post-Quantum Attacks
- 4. Solutions to Post-Quantum Attacks
- 5. Multi-Party Computation





• Cryptography (...), is the practice and study of techniques for secure communication in the presence of adversarial behavior".





• Cryptography (...), is the practice and study of techniques to ensure the confidentiality, integrity, authentication and non-repudiation of messages in the presence of adversarial behavior".

Confidentiality: keep sensitive information private.

Integrity: Assure the data received is valid.

Authentication: Prove who you are.

Non-repudiation: Emmiter can not deny sending a message.

### Security is a Belief System



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In Cryptography we work with hard problems.

A good example: inverting hash functions ( $h : \{0, 1\}^m \rightarrow \{0, 1\}^n$ ).

Pre-image resistance: Given an output y it is infisible to find any input x such that h(x) = y.

Second pre-image resistance: For a given specific input  $x_1$  it is infisible to find an input  $x_2$  such that  $h(x_1) = h(x_2)$ .

Collision resistance: It is hard to find any two inputs  $x_1$  and  $x_2$  such that  $h(x_1) = h(x_2)$ .

Let us toss a coin!

# **Symmetric Encryption**

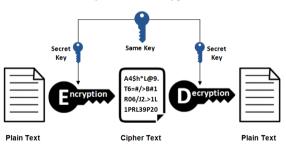


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Symmetric Encryption

- Emmiter and receiver have the same key.
- Fast with short signatures.



Symmetric Encryption

# **Asymmetric Encryption**

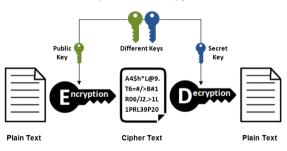


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Asymmetric Encryption

- Emmiter and receiver both have a key-pair.
- Slow with big signatures.



Asymmetric Encryption





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$$f: \{0, 1\}^m \to \{0, 1\}^n$$

- Brute-force approach:  $2^{n-1}$
- Grover's algorithm:  $2^{\frac{n}{2}}$
- Solution: Double key space  $(2^{\frac{2n}{2}})$ .

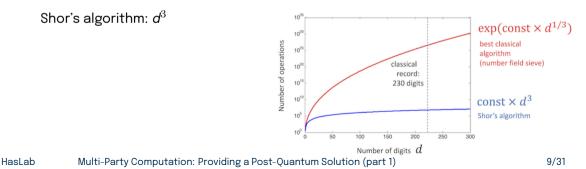




Factor an integer N with d decimal digits

Brute force approach:  $2^d$  (exponential in the number of digits d)

General Number Field Sieve approach:  $2^{\sqrt[3]{d}}$ 







In 2017, the American National Institute of Standards And Technologies (NIST) released a contest to find and standardize post-quantum secure schemes (NIST-PQC).

Key Encapsulation Mechanisms (KEM)

Digital Signature Schemes (DSS)

# **Key Encapsulation Mechanisms (KEM)**



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Chosen to Standardize:

CRYSTALS-Kyber (Lattice-based)

Postponed to the 4° round:

BIKE (Bit Flipping Key Encapsulation)

Classic McEliece (Binary Goppa Code)

HQC (Hamming Quasi-Cyclic)

SIKE (Supersingular Isogeny Key Encapsulation)

## **Digital Signature Schemes**



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Chosen to Standardize:

- CRYSTALS-Dilithium (Lattice-based)
- Falcon (Lattice-based)
- SPHINCS+ (hash-based)

All other schemes were dropped!

NIST is due to start a new competition to find and standardize new Digital Signature Schemes (NIST-PQC: DSS).



PICNIC (Multi-Party Computation) was also present in the 3° round.

- PICNIC was faster than SPHINCS+ but had larger signatures.
- "Obtaining further improvements under the same paradigm as Picnic (...) may (...) lead to a signature scheme with significantly better performance than the current design."
- In recent several other MPC schemes appeared with significant improments on PICNIC's signature size and speed.



Multiple parties computing a function's output jointly while keeping their respective shares private.

- A = 45
- $\mathsf{B}=65$
- C = 75

Avg(A,B,C) = 61.7



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Imagine we have three players:

Alice: 45 Bob: 65 Charlie: 75

Objective: Compute the average of their incomes without revealing those incomes.



First step: Each player divides its incomes into three shares.

Alice:  $45 \rightarrow (-11, 24, 32)$ Bob:  $65 \rightarrow (35, 45, -15)$ Charlie:  $75 \rightarrow (10, 15, 50)$ 

Second Step: Each player shares two of his/hers three shares.

Players	Alice	Bob	Charlie	Income
Alice	-11	24	32	45
Bob	35	45	-15	65
Charlie	10	15	50	75

### **Multi-Party Computation (MPC)**



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Third step: Sum each players' new shares.

Players	A	В	С	Income
Alice	-11	24	32	45
Bob	35	45	-15	65
Charlie	10	15	50	75
Computed Values	34	84	67	-

Fourth step: Compute the average of the new values.

Avg(45,65,75) = Avg(34,74,67) = 61,7

### **Multi-Party Computation in the head**



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Depending of the channel used latency can increase or decrease.

- Talking
- Email
- Letters

Solution: Do multi party computations in the head (Alice, Bob and Charlie become fictional).



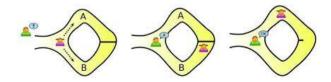


Zero-knowledge: Revealing we know a secret without disclosing it.

Bob want to go around the tunnel but does not know the pin code.

Alice knows the pin code and will sell it to Bob.

Bob wants to make sure Alice actually knows the secret code before paying.





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